Normalisation of closed linear $\lambda$-terms and patterns in trivalent maps


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Based on joint work with Olivier Bodini, Bernhard Gittenberger Michael Wallner, and Noam Zeilberger.

16th workshop on Computational Logic and Applications
Friday, January 13th 2023

The plan

- A brief overview of maps and the $\lambda$-calculus
- Context and results
- A strategy for deriving such results
- Normalisation of closed linear terms
- Other patterns in terms and maps

Quiz
asymptotically
Of the following types of redices, which one has the highest mean number of occurences in random closed linear terms?
a) Abstraction applied to variable: ( $\lambda x . t) y$
b) Abstraction applied to abstraction: ( $\lambda x . t$ ) ( $\lambda y . t^{\prime}$ )
c) Abstraction applied to application: ( $\lambda x . t$ ) (a b)

What are maps?


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$$
4 С Т \ldots
$$

- A central object in modern combinatorics, but not only that: probability, algebraic geometry, theoretical physics... scaling limits... matrix integrals, Witten's conjecture, ...

What are maps?


- A central object in modern combinatorics, but not only that: probability, algebraic geometry, theoretical physics...
- Their enumeration was pioneered by Cute in the 60 s , as part of his approach to the four colour theorem.

The untyped linear $\lambda$-calculus

- Its terms are formed inductively



## The untyped linear $\lambda$-calculus

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The untyped linear $\lambda$-calculus

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$$
\frac{\Gamma, x, y, \Delta \vdash \mathrm{t}}{\Gamma, y, x, \Delta \vdash \mathrm{t}} \text { exc } \quad \frac{\Gamma, x, y \vdash \mathrm{t}}{\Gamma, x \vdash \mathrm{t}[\mathrm{y}:=\mathrm{x}]} \text { con }
$$

## Linking terms and maps

- In 2013, Bodini, Gardy, Jacquot, describe a series of bijections:
rooted trivalent maps $\leftrightarrow$ closed linear terms
rooted (2,3)-valent maps $\leftrightarrow$ closed affine terms
In the same year, together with Gittenberger, they study:
$\operatorname{BCI}(p)$ terms (each bound variable appears $p$ times)
general closed $\lambda$-terms


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rooted planar maps $\leftrightarrow$ normal planar lambda terms
Both make use of decompositions in the style of Tutte! (cf. the approach of Arquès-Béraud in 2000)
- In 2015, Zeilberger advocates for
"linear lambda terms as invariants of rooted trivalent maps"

The thesis in context

Related work has been carried out on:

- Parameter studies on general $\lambda$-terms (ex., [BBD19]). different size notions!

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Our focus is on:

- Trivalent maps, linear terms, and related families.
- Exploring the combinatorial interplay of maps and terms.
- Study of pairs of parameters on maps and terms.


## Our results

Parameters on maps and terms of arbitrary genus (number of):

- Loops in trivalent maps and identity-subterms in closed linear terms

Limit law: Poisson(1)

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- Steps to reach normal form for closed linear terms

Asymptotic mean bound below by: $\frac{11 n}{240}$

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## Our strategy:

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- $T=u z+z L^{2}+z \partial_{\mathfrak{u}} L$

- $\mathrm{T}=u z+z^{2}+z \mathrm{~L}^{2}+2 z^{4} \partial_{z} \mathrm{~L}$
-10

- $\mathrm{T}=u z^{2}+z^{4}+z^{5} \frac{\partial}{\partial z}\left(\ln \left(\exp \left(z^{2} / 2\right) \odot \exp \left(z^{3} / 3+u z\right)\right)\right)$


Our strategy:

1) Track evolution of parameters in decompositions of maps $/ \lambda$-terms
different decompositions $\rightsquigarrow$ differential equations, Hadamard products, ...

2) Develop tools for rapidly growing coefficients, based on:

- Moment pumping
- Bender's theorem for compositions $\mathrm{F}(z, \mathrm{G}(z))$
- Coefficient asymptotics of Cauchy products

$$
\left[z^{n}\right](A(z) \cdot B(z)) \sim a_{n} b_{0}+a_{0} b_{n}+O\left(a_{n-1}+b_{n-1}\right)
$$

for $A, B, G$ divergent and $F$ analytic

From closed terms to maps


From closed terms to maps

$$
(\lambda x . x)(\lambda y \cdot(\lambda z . z y)(\lambda w \cdot \lambda u . w u))
$$



From closed terms to maps

$$
(\lambda x \cdot x)(\lambda y \cdot(\lambda z \cdot z y)(\lambda w \cdot \lambda u \cdot w u))
$$



Dictionary

- \# subterms $\leftrightarrow \#$ edges


From closed terms to maps
$(\lambda x . x)(\lambda y \cdot(\lambda z . z y)(\lambda w . \lambda u . w u))$


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- \# subterms $\leftrightarrow$ \# edges
- closed subterms $\leftrightarrow$ bridges
- using variables in order $\leftrightarrow$ planarity of maps

From closed terms to maps

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Dictionary

- \# subterms $\leftrightarrow$ \# edges
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Q: What if we erase the labels? Can we recover them?
A: Yes, via an exploration process! [BGJ13, BGGJ13,Z16]

Decomposing rooted trivalent maps

Decomposing rooted trivalent maps



Decomposing rooted trivalent maps


Decomposing rooted trivalent maps


Decomposing rooted trivalent maps and closed linear terms!
$\lambda x . x$

$$
(s t)
$$

$$
\lambda x . t=\lambda x . t \mid u:=(x u)] \text { or }
$$

$$
\rightarrow \lambda x .[[u:=(u x)]
$$

edges

$\downarrow$
$T(z)=z^{2}+z T(z)^{2}+$

subterms

## Computing with the $\lambda$-calculus

Dynamics of the $\lambda$-calculus: $\beta$-reductions

$$
\left(\left(\lambda x . t_{1}\right) t_{2}\right) \xrightarrow{\beta} t_{1}\left[x:=t_{2}\right]
$$

represents:

$$
\mathrm{f}=\chi \mapsto \mathrm{t}_{1}
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Examples of reductions

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\begin{aligned}
& ((\lambda x . x) y) \xrightarrow{\beta} x[x:=y]=y \\
& ((\lambda x .(\lambda y \cdot(y \quad y)) z))(a b)) \xrightarrow{\beta}(\lambda x .(z x))(a \quad \text { b }) \xrightarrow{\beta}(z(a b))
\end{aligned}
$$

A term with no redices is called a normal form

Computing with the $\lambda$-calculus
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$\beta$-normalisation terminates in deterministic number of steps
$\beta$-reduction in maps

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Normalisation of random closed linear terms

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$$
\begin{aligned}
& \text { Q: How many steps does it take for a random closed linear term } \\
& \text { to reach normal form? }
\end{aligned}
$$

A lower bound is given by the number of $\beta$-redices!
What is the number of $\beta$-redices in a random linear $\lambda$-term?

uniform distribution

Normalisation of random closed linear terms $\rightarrow$ well defined! (strong normalisation + diamond)
Q: How many steps does it take for a random closed linear term to reach normal form?

A lower bound is given by the number of $\beta$-redices!
What is the number of $\beta$-redices in a random linear $\lambda$-term?

uniform distribution

Q: Why is this just a lower bound?
A: Because reducing a redex can create a new one!

## Mean number of $\beta$-redices in closed terms

 Tracking redices: starts off easy...Mean number of $\beta$-redices in closed terms Tracking redices: starts off easy...
loops
9

Mean number of $\beta$-redices in closed terms
Tracking redices: starts off easy...


Mean number of $\beta$-redices in closed terms
Tracking redices: then gets harder!
Abstractions, subcase 1.1


Mean number of $\beta$-redices in closed terms
Translating to a differential equation and pumping

$$
\begin{aligned}
\mathrm{T}=z^{2} & \left.+z \mathrm{~T}^{2}+z^{3}(1+(\mathrm{r}-1) z \mathrm{~T})\left(\frac{z(\mathrm{r}+5) \partial_{z} \mathrm{~T}}{3}-\left(\mathrm{r}^{2}-1\right) \partial_{\mathrm{r}} \mathrm{~T}\right)\right) \\
& +\frac{z^{4}(\mathrm{r}-1)^{2} \mathrm{~T}^{2}}{3}+\frac{4 z^{3}(\mathrm{r}-1) \mathrm{T}}{3}
\end{aligned}
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\end{aligned}
$$

Let $X_{n}$ be the random variable given by number of redices in a closed linear term of size $n \in 3 \mathbb{N}+2$. Then

$$
\begin{aligned}
& \mathbb{E}\left(X_{n}\right) \sim \frac{n}{24} \\
& \mathbb{V}\left(X_{n}\right) \sim \frac{n}{24}
\end{aligned}
$$

A lower bound for normalisation
Refining our counting to track reproducing redices:

A lower bound for normalisation
(see JJ Lévy's thesis)

Refining our counting to track reproducing redices:


$$
\begin{aligned}
& p_{1}=(\lambda x \cdot C[(x u)])(\lambda y \cdot t) \xrightarrow{\beta} C[((\lambda y \cdot t) u)] \\
& p_{2}=(\lambda x \cdot x)\left(\lambda y \cdot t_{1}\right) t_{2} \xrightarrow{\beta}\left(\lambda y \cdot t_{1}\right) t_{2} \\
& p_{3}=\left(\left(\lambda x \cdot \lambda y \cdot t_{1}\right) t_{2}\right) t_{3} \xrightarrow{\beta}\left(\lambda y \cdot t_{1}\left[x:=t_{2}\right]\right) t_{3}
\end{aligned}
$$

## Enumerating $p_{1}$-patterns

- Tracking the creation/destruction of patterns during the recursive decomposition:

Enumerating $p_{1}$-patterns

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Cuts destroying a $p_{1}$-pattern:


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- Tracking the creation/destruction of patterns during the recursive decomposition:

Cuts creating a $p_{1}$-pattern:


Thus we also need to keep track of:

$$
\left.C_{1}\left[\lambda x . C_{2}\left[\left(t_{1} x\right)\right]\right)\left(\lambda y . t_{2}\right)\right] \quad C_{1}\left[(\lambda x . x)\left(\lambda y . t_{2}\right)\right]
$$

## Enumerating $p_{1}$-patterns

- Tracking the creation/destruction of patterns during the recursive decomposition:

Applications creating $p_{1}$ and auxilliary patterns:


Thus, for an app. of the form ( $l_{1} \lambda y . t_{1}$ ) we need to consider how $l_{1}$ was formed.

## Enumerating $p_{1}$-patterns

-Thus we have the following equations:

$$
S=\Lambda+A
$$

$$
\Lambda=z^{2}+2 z^{4} S_{z}+(v-u+4(1-u)) z^{3} S_{u}+(u-v+4(1-v)) z^{3} S_{v}
$$

$$
A=z S^{2}+(u-1) z\left(z^{4} S_{z}+(v-u+2(1-u)) z^{3} S_{u}+2(1-v) z^{3} S_{v}\right) \cdot \Lambda
$$

$$
+(v-1) z\left(z^{2}+z^{4} S_{z}+(u-v+2(1-v)) z^{3} \mathrm{~S}_{\mathfrak{u}}+2(1-u) z^{3} \mathrm{~S}_{\mathfrak{u}}\right) \cdot \Lambda
$$

- Extracting the mean:
$\left.\partial_{\mathcal{u}} S\right|_{\mathcal{u}=1, v=1}$
$=\left.\left(2 z S \partial_{\mathfrak{u}} S+2 z^{4} \partial_{z, u} S+z^{7} \partial_{z} S+2 z^{9}\left(\partial_{z} S\right)^{2}-5 z^{3} \partial_{\mathfrak{u}} S+z^{3} \partial_{v} S\right)\right|_{\mathfrak{u}=1, v=1}$


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$$
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A= & z S^{2}+(u-1) z\left(z^{4} S_{z}+(v-u+2(1-u)) z^{3} S_{u}+2(1-v) z^{3} S_{v}\right) \cdot \Lambda \\
& +(v-1) z\left(z^{2}+z^{4} S_{z}+(u-v+2(1-v)) z^{3} \mathbf{S}_{u}+2(1-u) z^{3} S_{u}\right) \cdot \Lambda
\end{aligned}
$$

- Extracting the mean:



$$
\begin{aligned}
& \left.\partial_{\mathcal{u}} S\right|_{u=1, v=1} \\
& \text { bijection: } \partial_{v} \leftrightarrow \partial_{u} \\
& =\left.\left(2 z S \partial_{u} S+2 z^{4} \partial_{z, u} S+z^{7} \partial_{z} S+2 z^{9}\left(\partial_{z} S\right)^{2}-5 z^{3} \partial_{u} S+z^{3} \partial_{v} S\right)\right|_{u=1, v=1}
\end{aligned}
$$

## Enumerating $p_{1}$-patterns

- Finally we obtain a mean number of occurences:

$$
\mathbb{E}\left[\# p_{1} \text { patterns }\right] \sim \frac{1}{6}
$$

Enumerating $p_{1}$-patterns and $p_{2}$-patterns

- Finally we obtain a mean number of occurences:

$$
\mathbb{E}\left[\# p_{1} \text { patterns }\right] \sim \frac{1}{6}
$$

- Analogously, we have a mean number of occurences for $\mathrm{p}_{2}$ :

$$
\mathbb{E}\left[\# p_{2} \text { patterns }\right] \sim \frac{1}{48}
$$

Both are asymptotically constant in expectation!

## Enumerating $p_{3}$-patterns

- As before, we'll also need to enumerate auxilliary patterns:

$$
\left(\lambda x . \lambda y . t_{1}\right) \quad\left(\lambda x . \lambda y . t_{1}\right) t_{2} \quad\left(\lambda x . \lambda y . t_{1}\right) t_{2} t_{3}
$$

- However we run into a problem:



## Enumerating $p_{3}$-patterns

- Generatingfunctionology fails, we revert to more elementary methods:

$$
\mathbb{E}\left(\mathrm{V}_{n}\right)=\mathbb{E}\left(\mathrm{V}_{n} \mid \Lambda_{n}\right) \cdot \frac{\left|\Lambda_{n}\right|}{\left|\mathrm{L}_{n}\right|}+\mathbb{E}\left(\mathrm{V}_{n} \mid A_{n}\right) \cdot \frac{\left|A_{n}\right|}{\left|\mathrm{L}_{n}\right|}
$$

## Enumerating $p_{3}$-patterns

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$$

Magic: linear over families of all possible abstractions created via cuts from a fixed term!

$$
\begin{aligned}
& \bar{X}_{n}=(2 n-12) \bar{X}_{n-3} 2 \bar{Y}_{n-3} \\
& \bar{Y}_{n}=(2 n-6) Y_{n-3}-6 Y_{n-3} \\
& \bar{Z}_{n}=2(n-4)\left(Z+\mathbf{1}_{\wedge_{n}}\right)
\end{aligned}
$$

where: $X_{n}$ counts $\#$ of $p_{1}$ patt. over terms of size $n$ $Y_{n}$ is the same for the pattern $\left(\lambda x . \lambda y . t_{1}\right) t_{2}$, and $Z$ is the same for the pattern ( $\lambda x . \lambda y . t_{1}$ )
The $\bar{V}$ for $V \in\left\{X_{n}, Y_{n}, Z_{n}\right\}$ are cummulatives over families of abstractions

The lower bound

Theorem
Let $W_{n}$ be the random variable given by number of steps required to normalise a linear term of size $n \in 3 \mathbb{N}+2$. Then

$$
\mathbb{E}\left(W_{n}\right) \geqslant \frac{11 n}{240}, \text { for } n \text { large enough }
$$

During an open problem session of CLA 2020, Noam Zeilberger conjectured:

$$
\mathbb{E}\left(W_{n}\right) \sim \frac{n}{21}
$$

We got pretty close:

$$
\frac{11}{240}-\frac{1}{21}=0.001 \ldots
$$

Counting redices by type of argument


Counting redices by type of argument


## Theorem

$\mathbb{E}(\#(\lambda x . t) y) \sim \frac{n}{30}$
$\mathbb{E}\left(\#(\lambda x . t)\left(\lambda y . t^{\prime}\right)\right) \sim \frac{1}{20}$
$\mathbb{E}\left(\#(\mathrm{ab})\left(\lambda y \cdot \mathrm{t}^{\prime}\right)\right) \sim \frac{n}{120}$

Open problems

- Classify patterns according to their expected number of occurences: constant or linear in $n$.

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- Automatise the process of obtaining specifications tracking occurences of our desired patterns.

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- Explore the meaning of $\beta$-reduction om maps. Connections to knot theory?

Open problems

- Classify patterns according to their expected number of occurences: constant or linear in $n$.
- Automatise the process of obtaining specifications tracking occurences of our desired patterns.
- Explore the meaning of $\beta$-reduction om maps. Connections to knot theory?

Thank you!

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[^0]:    well defined! (strong normalisation + diamond)
    Q: How many steps does it take for a random closed linear term to reach normal form?

